# A recursive construction for simple $t$-designs using resolutions 

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#### Abstract

This work presents a recursive construction for simple $t$-designs using resolutions of the ingredient designs. The result extends a construction of $t$-designs in our recent paper [39]. Essentially, the method in [39] describes the blocks of a constructed design as a collection of block unions from a number of appropriate pairs of disjoint ingredient designs. Now, if some pairs of these ingredient $t$-designs have both suitable $s$-resolutions, then we can define a distance mapping on their resolution classes. Using this mapping enables us to have more possibilities for forming blocks from those pairs. The method makes it possible for constructing many new simple $t$-designs. We give some application results of the new construction.


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## 1 Introduction

In a recent paper [39] we have presented a recursive method for constructing simple $t$-designs for arbitrary $t$. The method is of combinatorial nature since it requires finding solutions for the indices of the ingredient designs that satisfy a certain set of equalities. In essence, the core of the construction is that the blocks of a constructed design are built as a collection of block unions from a number of appropriate pairs of disjoint ingredient designs. In particular, when a pair of ingredient designs is used, we take as new blocks the unions of all the pairs of blocks in the two ingredient designs. For the sake of simplicity we refer to this construction method as the basic method or the basic construction.

In the present paper we describe an extension of the basic construction by assuming that a subset of pairs of ingredient designs have suitable resolutions. For those given pairs we may define a distance mapping on their resolution classes. By
using this mapping we have more possibilities for forming blocks from those pairs other than taking the unions of all possible pairs of blocks in the ingredient designs. This construction actually extends the basic construction since many new simple $t$ designs can only be constructed with the new method. The crucial point of this extension is the use of $s$-resolutions for $t$-designs. The concept of $s$-resolutions may be viewed as a generalization of the notion of parallelisms, which may be termed as $(1,1)$-resolutions, i.e. the blocks of the $t$-design can be partitioned into classes of mutually disjoint blocks such that every point is in exactly one block of each class. To date very little is known about $s$-resolutions for $t$-designs when $s \geq 2$, except for the trivial $t$-designs. In this case, an $s$-resolution of the trivial $t$-design turns out to be a large set of $s$-designs. A great deal of results about large sets of $s$-designs have been achieved by many researchers, see the references below. We will describe our construction in terms of $s$-resolutions for $t$-designs in general. However we will restrict its applications just for the case where pairs of trivial designs are used and each has a suitable large set. Even with this limitation we find that the construction using resolutions still possesses its strength since it produces many simple $t$-designs.

It is worthwhile to emphasize that constructing simple $t$-designs for large $t$ is a challenging problem in design theory. There are several major approaches to the problem. These include constructing $t$-designs from large sets of $t$-designs, for instance $[1,18,13,16,19,21,23,24,25,32,33,34,41]$; constructing $t$-designs by using prescribed automorphism groups, for example $[2,3,6,7,8,9,10,14,20,22,26,29]$; or contructing $t$-designs via recursive construction methods, see for instance [15, 17, 27, 31, 30, 36, 37, 38, 39, 40].

## 2 Preliminaries

We recall some basic definitions. A $t$-design, denoted by $t-(v, k, \lambda)$, is a pair $(X, \mathfrak{B})$, where $X$ is a $v$-set of points and $\mathfrak{B}$ is a collection of $k$-subsets, called blocks, of $X$ having the property that every $t$-set of $X$ is a subset of exactly $\lambda$ blocks in $\mathfrak{B}$. The parameter $\lambda$ is called the index of the design. A $t$-design is called simple if no two blocks are identical i.e. no block of $\mathcal{B}$ is repeated; otherwise, it is called non-simple (i.e. $\mathfrak{B}$ is a multiset). It can be shown by simple counting that a $t-(v, k, \lambda)$ design is an $s-\left(v, k, \lambda_{s}\right)$ design for $0 \leq s \leq t$, where $\lambda_{s}=\lambda\binom{v-s}{t-s} /\binom{k-s}{t-s}$. Since $\lambda_{s}$ is an integer, necessary conditions for the parameters of a $t$-design are $\binom{k-s}{t-s} \left\lvert\, \lambda\binom{v-s}{t-s}\right.$, for $0 \leq s \leq t$. For given $t, v$ and $k$, we denote by $\lambda_{\min }(t, k, v)$, or $\lambda_{\min }$ for short, the smallest positive integer such that these conditions are satisfied for all $0 \leq s \leq t$. By complementing each block in $X$ of a $t-(v, k, \lambda)$ design, we obtain a $t-\left(v, v-k, \lambda^{*}\right)$ design with $\lambda^{*}=\lambda\binom{v-k}{t} /\binom{k}{t}$, hence we shall assume that $k \leq v / 2$. The largest value for $\lambda$ for which a simple $t-(v, k, \lambda)$ design exists is denoted by $\lambda_{\max }$ and we have $\lambda_{\max }=\binom{v-t}{k-t}$. The simple $t-\left(v, k, \lambda_{\max }\right)$ design is called the complete design or the trivial design. A $t$ - $(v, k, 1)$ design is called a $t$-Steiner system.

We refer the reader to $[5,12]$ for more information about designs.
Definition 2.1 $A t-(v, k, \lambda)$-design $(X, \mathfrak{B})$ is said to be $(s, \tau)$-resolvable with $0<s<$ $t$, if its block set $\mathfrak{B}$ can be partitioned into $N$ classes $\mathfrak{A}_{1}, \ldots, \mathfrak{A}_{N}$ such that $\left(X, \mathfrak{A}_{i}\right)$ is
an $s-(v, k, \tau)$ design for all $i=1, \ldots, N$. Each $\mathfrak{A}_{i}$ is called a resolution class. We also say that a $t-(v, k, \lambda)$-design has an $s$-resolution, if it is $(s, \tau)$-resolvable for a certain $\tau$.

It is worth noting that the concept of resolvability (i.e. ( 1,1 )-resolvability) for BIBD introduced by Bose in 1942 [11] was generalized by Shrikhande and Raghavarao to $\tau$-resolvability (i.e. $(1, \tau)$-resolvability) for BIBD in 1963 [28]. A definition of $(s, \lambda)$ resolvability for $t$-designs with $t \geq 3$ may be found in [4]. In that paper Baker shows that the Steiner quadruple system $3-\left(4^{m}, 4,1\right)$ constructed from an even dimensional affine space over the field of two elements has a $(2,1)$-resolution. Also, Teirlinck shows for example that there exists a 2 -resolvable $3-\left(2 p^{n}+2,4,1\right)$ design with $p \in$ $\{7,31,127\}$, for any positive integer $n$, [35].

To date, very little is known about $s$-resolutions of non-trivial $t-(v, k, \lambda)$ designs for $t \geq 3$ and $s \geq 2$. Here are examples with $t=4$ and $s=3$. In [2], Alltop has shown that there exists a simple $4-(q+1,5,5)$ design for every $q=2^{n}, n \geq 5, n$ odd. This 4 -design $(X, \mathfrak{B})$ is constructed by using the group $\operatorname{PGL}(2, q)$, which acts sharply 3-transitively on the projective line $X=\operatorname{GF}(q) \cup\{\infty\}$. The block set $\mathfrak{B}$ is a disjoint union of $(q-2) / 6$ orbits of 5 -sets of $X$ under PGL $(2, q)$. Each orbit forms a $3-(q+1,5,15)$ design. Hence each $4-(q+1,5,5)$ design in the Alltop's family has a $(3,15)$-resolution.

When $(X, \mathfrak{B})$ is the trivial $t-\left(v, k,\binom{v-t}{k-t}\right)$ design, then an $(s, \tau)$-resolution of $(X, \mathfrak{B})$ is called a large set. Thus, a large set is a partition of the complete $t-\left(v, k,\binom{v-t}{k-t}\right)$ design into $s-(v, k, \tau)$ designs, and is denoted by $\operatorname{LS}[N](s, k, v)$, where $N=\binom{v-s}{k-s} / \tau$ is the number of resolution classes in the partition.

We define a distance on the resolution classes of a $t$-design as follows.
Definition 2.2 Let $D$ be at- $(v, k, \lambda)$ design admitting an $(s, \tau)$-resolution with $\mathfrak{A}_{1}, \ldots, \mathfrak{A}_{N}$ as resolution classes. Define a distance between any two classes $\mathfrak{A}_{i}$ and $\mathfrak{A}_{j}$ by $d\left(\mathfrak{A}_{i}, \mathfrak{A}_{j}\right)=$ $\min \{|i-j|, N-|i-j|\}$.

### 2.1 The basic construction

In this section, we summarize the basic construction as described in [39]. This preparation is necessary for the description of the construction using resolution in the next section.

We first give notation and definitions. Let $t, v, k$ be non-negative integers such that $v \geq k \geq t \geq 0$. Let $X$ be a $v$-set and let $X=X_{1} \cup X_{2}$ be a partition of $X$ (i.e $X_{1} \cap X_{2}=\emptyset$ ) with $\left|X_{1}\right|=v_{1}$ and $\left|X_{2}\right|=v_{2}$.

The parameter set $t-\left(v_{2}, j, \bar{\lambda}_{t}^{(j)}\right)$ for a design indicates that the point set of the design is $X_{2}$. Also, a design defined on the point set $X_{2}$ is denoted by $\bar{D}=\left(X_{2}, \overline{\mathfrak{B}}\right)$.
(i) For $i=0, \ldots, t$, let $D_{i}=\left(X_{1}, \mathfrak{B}^{(i)}\right)$ be the complete $i-\left(v_{1}, i, 1\right)$ design. For $i=t+1, \ldots, k$, let $D_{i}=\left(X_{1}, \mathfrak{B}^{(i)}\right)$ be a simple $t-\left(v_{1}, i, \lambda_{t}^{(i)}\right)$ design.
(ii) Similarly, for $i=0, \ldots, t$, let $\bar{D}_{i}=\left(X_{2}, \overline{\mathfrak{B}}^{(i)}\right)$ be the complete $i$ - $\left(v_{2}, i, 1\right)$ design. And for $i=t+1, \ldots, k$, let $\bar{D}_{i}=\left(X_{2}, \overline{\mathfrak{B}}^{(i)}\right)$ be a simple $t-\left(v_{2}, i, \bar{\lambda}_{t}^{(i)}\right)$ design.
(iii) Two degenerate cases for designs occur when either $k=t=0$ or $v=k$. The first case $k=t=0$ gives an "empty" design, denoted by $\emptyset$, however we use the convention that the number of blocks of the empty design is 1 (i.e. the unique block is the empty block). The second case $v=k$ gives a degenerate $k$-design having just 1 block consisting of all $v$ points. Thus, in these two extreme cases the number of blocks of the designs is always 1 .
(iv) We denote by $T_{(r, t-r)}$ a $t$-subset $T$ of $X$ with $\left|T \cap X_{1}\right|=r$ and hence $\left|T \cap X_{2}\right|=$ $t-r$, for $r=0, \ldots, t$. It is clear that any $t$-subset of $X$ is a $T_{(r, t-r)}$ set for some $r \in\{0, \ldots, t\}$.
(v) Let $X$ be a finite set and let $u \in\{0,1\}$. The notation $X \times[u]$ has the following meaning. $X \times[0]$ is the empty set $\emptyset$, and $X \times[1]=X$.

The basic construction in [39] is as follows.
Consider $(k+1)$ pairs of simple designs $\left(D_{i}, \bar{D}_{k-i}\right)$ for $i=0, \ldots, k$, where $D_{i}=$ $\left(X_{1}, \mathfrak{B}^{(i)}\right)$ is a simple $t-\left(v_{1}, i, \lambda_{t}^{(i)}\right)$ design and $\bar{D}_{k-i}=\left(X_{2}, \overline{\mathfrak{B}}^{(k-i)}\right)$ a simple $t-\left(v_{2}, k-\right.$ $\left.i, \bar{\lambda}_{t}^{(k-i)}\right)$ design, as defined above. For each pair $\left(D_{i}, \bar{D}_{k-i}\right)$ define

$$
\mathfrak{B}_{(i, k-i)}:=\left\{B=B_{i} \cup \bar{B}_{k-i} \mid B_{i} \in \mathfrak{B}^{(i)}, \bar{B}_{k-i} \in \overline{\mathfrak{B}}^{(k-i)}\right\} .
$$

Define

$$
\mathfrak{B}:=\mathfrak{B}_{(0, k)} \times\left[u_{0}\right] \cup \mathfrak{B}_{(1, k-1)} \times\left[u_{1}\right] \cup \cdots \cup \mathfrak{B}_{(k-1,1)} \times\left[u_{k-1}\right] \cup \mathfrak{B}_{(k, 0)} \times\left[u_{k}\right],
$$

where $u_{i} \in\{0,1\}$, for $i=0, \ldots, k$.
It should be remarked that the notation $\mathfrak{B}_{(i, k-i)} \times\left[u_{i}\right]$, as defined in (v) above, indicates that either we have an empty set $\emptyset$ (when $u_{i}=0$ ) or the set $\mathfrak{B}_{(i, k-i)}$ itself (when $u_{i}=1$ ). The empty set case means that the pair $\left(D_{i}, \bar{D}_{k-i}\right)$ is not used and the other case means the pair $\left(D_{i}, \bar{D}_{k-i}\right)$ is used.

It can be shown that for a given $t$-set $T_{(r, t-r)}$ of $X$ the number of blocks in $\mathfrak{B}$ containing $T_{(r, t-r)}$ is equal to

$$
L_{r, t-r}:=\sum_{i=0}^{k} u_{i} \cdot \lambda_{r}^{(i)} \cdot \bar{\lambda}_{t-r}^{(k-i)}
$$

Therefore, if

$$
L_{0, t}=L_{1, t}=L_{2, t-2}=\cdots=L_{t, 0}:=\Lambda,
$$

where $\Lambda$ is a positive integer, then $(X, \mathfrak{B})$ forms a simple $t$-design with parameters $t-(v, k, \Lambda)$.

We record the basic construction in the following theorem.

Theorem 2.1 (Basic construction) Let $v, k, t$ be integers with $v>k>t \geq 2$. Let $X$ be a v-set and let $X=X_{1} \cup X_{2}$ be a partition of $X$ with $\left|X_{1}\right|=v_{1}$ and $\left|X_{2}\right|=v_{2}$. Let $D_{i}=\left(X_{1}, \mathfrak{B}^{(i)}\right)$ be the complete $i-\left(v_{1}, i, 1\right)$ design for $i=0, \ldots, t$ and let $D_{i}=\left(X_{1}, \mathfrak{B}^{(i)}\right)$ be a simple $t-\left(v_{1}, i, \lambda_{t}^{(i)}\right)$ design for $i=t+1, \ldots, k$. Similarly, let $\bar{D}_{i}=\left(X_{2}, \overline{\mathfrak{B}}^{(i)}\right)$ be the complete $i$ - $\left(v_{2}, i, 1\right)$ design for $i=0, \ldots, t$, and let $\bar{D}_{i}=$ $\left(X_{2}, \overline{\mathfrak{B}}^{(i)}\right)$ be a simple $t-\left(v_{2}, i, \bar{\lambda}_{t}^{(i)}\right)$ design for $i=t+1, \ldots, k$. Define

$$
\mathfrak{B}=\mathfrak{B}_{(0, k)} \times\left[u_{0}\right] \cup \mathfrak{B}_{(1, k-1)} \times\left[u_{1}\right] \cup \cdots \cup \mathfrak{B}_{(k-1,1)} \times\left[u_{k-1}\right] \cup \mathfrak{B}_{(k, 0)} \times\left[u_{k}\right],
$$

where

$$
\mathfrak{B}_{(i, k-i)}=\left\{B=B_{i} \cup \bar{B}_{k-i} \mid B_{i} \in \mathfrak{B}^{(i)}, \bar{B}_{k-i} \in \overline{\mathfrak{B}}^{(k-i)}\right\}
$$

Assume that

$$
\begin{equation*}
L_{0, t}=L_{1, t-1}=L_{2, t-2}=\cdots=L_{t, 0}:=\Lambda \tag{1}
\end{equation*}
$$

for a positive integer $\Lambda$, where

$$
\begin{equation*}
L_{r, t-r}=\sum_{i=0}^{k} u_{i} \cdot \lambda_{r}^{(i)} \cdot \bar{\lambda}_{t-r}^{(k-i)} \tag{2}
\end{equation*}
$$

$r=0, \ldots, t$, and $u_{i} \in\{0,1\}$, for $i=0, \ldots, k$. Then $(X, \mathfrak{B})$ is a simple $t-(v, k, \Lambda)$ design.

## 3 The construction using resolutions

In this section we describe a recursive construction of simple $t$-designs using resolutions. Note that in the basic construction, if a pair ( $D_{i}, \bar{D}_{k-i}$ ) is used in the construction (i.e. $u_{i}=1$ ), then the new blocks formed by this pair consist of taking the union of each block of $D_{i}$ with each block of $\bar{D}_{k-i}$. The crucial idea of the construction using resolutions is that if $D_{i}$ and $\bar{D}_{k-i}$ have appropriate $s_{1^{-}}$and $s_{2}$-resolutions with the same number of resolution classes, then the new blocks are formed according to the distance mapping defined on the resolution classes of $D_{i}$ and $\bar{D}_{k-i}$ rather than taking the unions of each block of $D_{i}$ with each block of $\bar{D}_{k-i}$.

In the following we go into detail of the construction. We make use of the notation and definitions for the basic construction in the previous section. When for a certain $i \in\{0, \ldots, k\}$ the $t$ - $\left(v_{1}, i, \lambda_{t}^{(i)}\right)$ design $D_{i}=\left(X_{1}, \mathfrak{B}^{(i)}\right)$ has an $s_{i}$-resolution, i.e. $D_{i}$ can be partitioned into $N_{i}$ disjoint $\left(X_{1}, \mathfrak{A}_{h}^{(i)}\right)$ designs with parameters $s_{i}-\left(v_{1}, i, \lambda_{s_{i}}^{*(i)}\right)$, $s_{i}<t$, then we write

$$
\mathfrak{B}^{(i)}=\bigcup_{h=1}^{N_{i}} \mathfrak{A}_{h}^{(i)},
$$

where

$$
N_{i}=\lambda_{t}^{(i)}\binom{v_{1}-s_{i}}{t-s_{i}} / \lambda_{s_{i}}^{*(i)}\binom{i-s_{i}}{t-s_{i}}
$$

Similarly, we write

$$
\overline{\mathfrak{B}}^{(k-i)}=\bigcup_{h=1}^{\bar{N}_{k-i}} \overline{\mathfrak{A}}_{h}^{(i)},
$$

when the blocks of a $t-\left(v_{2}, k-i, \bar{\lambda}_{t}^{(k-i)}\right)$ design $\bar{D}_{k-i}=\left(X_{2}, \overline{\mathfrak{B}}^{(k-i)}\right)$ can be partitioned into $\bar{N}_{k-i}$ disjoint $\left(X_{2}, \overline{\mathfrak{A}}_{h}^{(k-i)}\right)$ designs with parameters $s_{k-i}\left(v_{2}, k-i, \bar{\lambda}_{s_{k-i}}^{*(k-i)}\right)$, where

$$
\bar{N}_{k-i}=\bar{\lambda}_{t}^{(k-i)}\binom{v_{2}-s_{k-i}}{t-s_{k-i}} / \bar{\lambda}_{s_{k-i}}^{*(k-i)}\binom{k-i-s_{k-i}}{t-s_{k-i}}
$$

is the number of $s_{k-i}$-resolution classes.
Let $K=\{(0, k),(1, k-1), \ldots,(k-1,1),(k, 0)\}$. Assume there exists a subset $R \subseteq K$ such that if $(i, k-i) \in R$, then $D_{i}$ and $\bar{D}_{k-i}$ have an $s_{i}$-resolution of size $N_{i}$ and an $s_{k-i}$-resolution of size $\bar{N}_{k-i}$, respectively, satisfying the following conditions.
(i) $N_{i}=\bar{N}_{k-i}$,
(ii) $s_{i}+s_{k-i} \geq 2\left\lfloor\frac{t}{2}\right\rfloor$.

The construction consists of building two types of blocks.
(1) For each pair $(i, k-i) \in K \backslash R$ form a subset of new blocks $\mathfrak{B}_{(i, k-i)}$ from the pair $\left(D_{i}, \bar{D}_{k-i}\right)$ as

$$
\mathfrak{B}_{(i, k-i)}:=\left\{B=B_{i} \cup \bar{B}_{k-i} \mid B_{i} \in \mathfrak{B}^{(i)}, \bar{B}_{k-i} \in \overline{\mathfrak{B}}^{(k-i)}\right\} .
$$

(2) For each pair $(i, k-i) \in R$ form a subset of new blocks $\mathfrak{B}_{(i, k-i)}^{*}$ from $\left(D_{i}, \bar{D}_{k-i}\right)$ by using an $s_{i}$-resolution of $D_{i}$ and an $s_{k-i}$-resolution of $\bar{D}_{k-i}$ as follows.
$\mathfrak{B}_{(i, k-i)}^{*}:=\left\{B_{i} \cup \bar{B}_{k-i} \mid B_{i} \in \mathfrak{A}_{h}^{(i)}, \bar{B}_{k-i} \in \overline{\mathfrak{A}}_{j}^{(k-i)}, \varepsilon_{i} \leq d\left(\mathfrak{A}_{h}^{(i)}, \mathfrak{A}_{j}^{(i)}\right) \leq w_{i}, \varepsilon_{i}=0,1 ; w_{i} \leq\left\lfloor\frac{N_{i}}{2}\right\rfloor\right\}$.
Further, define

$$
z_{i}:=\left(2 w_{i}+1-\varepsilon_{i}\right), \text { if } w_{i}<\frac{N_{i}}{2}, \text { and } z_{i}:=\left(2 w_{i}-\varepsilon_{i}\right), \text { if } w_{i}=\frac{N_{i}}{2} .
$$

Note that $w_{i}$ and $z_{i}$ are considered as variables.
Now, let $T_{(r, t-r)}$ be a $t$-set of $X$ for $r=0, \ldots, t$. According to the property of $s_{i}$ and $s_{k-i}$ one of the following cases has to occur.
(a) $r \leq s_{i}$ and $t-r \leq s_{k-i}$. Then $T_{(r, t-r)}$ is contained in

$$
\Lambda_{r, t-r}^{*(i, k-i)}=\lambda_{r}^{*(i)} \cdot \lambda_{t-r}^{*(k-i)} \cdot N_{i} . z_{i}
$$

blocks of $\mathfrak{B}_{(i, k-i)}^{*}$.
(b) $r \leq s_{i}$ and $t-r>s_{k-i}$. Then $T_{(r, t-r)}$ is contained in

$$
\Lambda_{r, t-r}^{*(i, k-i)}=\lambda_{r}^{*(i)} \cdot \bar{\lambda}_{t-r}^{(k-i)} \cdot z_{i}
$$

blocks of $\mathfrak{B}_{(i, k-i)}^{*}$.
(c) $r>s_{i}$ and $t-r \leq s_{k-i}$. Then $T_{(r, t-r)}$ is contained in

$$
\Lambda_{r, t-r}^{*(i, k-i)}=\lambda_{r}^{(i)} \cdot \bar{\lambda}_{t-r}^{*(k-i)} \cdot z_{i}
$$

blocks of $\mathfrak{B}_{(i, k-i)}^{*}$.

It is straightforward to verify the values of $\Lambda_{r, t-r}^{*(i, k-i)}$ for the cases (a), (b) and (c) above. In case (a) each $r$-subset of $X_{1}$ is contained in $\lambda_{r}^{*(i)}$ blocks of $\mathfrak{A}_{h}^{(i)}$ and each $(t-r)$-subset of $X_{2}$ in $\bar{\lambda}_{t-r}^{*(k-i)}$ blocks of $\overline{\mathfrak{A}}_{j}^{(k-i)}$. Thus each pair $\left(\mathfrak{A}_{h}^{(i)}, \overline{\mathfrak{A}}_{j}^{(k-i)}\right)$ contributes $\lambda_{r}^{*(i)} \cdot \lambda_{t-r}^{*(k-i)}$ blocks to $\mathfrak{B}_{(i, k-i)}^{*}$. Now each of the $N_{i}$ resolution classes $\mathfrak{A}_{1}^{(i)}, \ldots, \mathfrak{A}_{N_{i}}^{(i)}$ is combined with $z_{i}$ resolution classes of $\overline{\mathfrak{A}}_{1}^{(k-i)}, \ldots, \overline{\mathfrak{A}}_{N_{i}}^{(k-i)}$, therefore we have $\Lambda_{r, t-r}^{*(i, k-i)}=\lambda_{r}^{*(i)} \cdot \bar{\lambda}_{t-r}^{*(k-i)} \cdot N_{i} \cdot z_{i}$

In case (b) each $r$-subset of $X_{1}$ is contained in $\lambda_{r}^{*(i)}$ blocks of $\mathfrak{A}_{h}^{(i)}$ and each $(t-r)$ subset of $X_{2}$ in $\bar{\lambda}_{t-r}^{(k-i)}$ blocks of $\overline{\mathfrak{B}}^{(k-i)}$. These blocks are distributed in the $N_{i}$ resolution classes $\overline{\mathfrak{A}}_{1}^{(k-i)}, \ldots, \overline{\mathfrak{A}}_{N_{i}}^{(k-i)}$. Each class $\overline{\mathfrak{A}}_{j}^{(k-i)}$ is combined $z_{i}$ times with $\mathfrak{A}_{h}^{(i)}$. Hence, in this case, the contribution of the blocks to $\mathfrak{B}_{(i, k-i)}^{*}$ is $\Lambda_{r, t-r}^{*(i, k-i)}=\lambda_{r}^{*(i)} \cdot \bar{\lambda}_{t-r}^{(k-i)} . z_{i}$.

The case (c) is similar to case (b).
Define

$$
\mathfrak{B}:=\bigcup_{(i, k-i) \in R} \mathfrak{B}_{(i, k-i)}^{*} \times\left[u_{i}\right] \cup \bigcup_{(i, k-i) \in K \backslash R} \mathfrak{B}_{(i, k-i)} \times\left[u_{i}\right],
$$

with $u_{i} \in\{0,1\}, i=0, \ldots, k$.
The above presentation can be summarized as follows. Let $T_{(r, t-r)}$ be a $t$-subset of $X$ for $r=0, \ldots, t$. The number of blocks in $\mathfrak{B}_{(i, k-i)}$ containing $T_{(r, t-r)}$, for all $(i, k-i) \in K \backslash R$, is then

$$
\sum_{(i, k-i) \in K \backslash R} u_{i} \cdot \Lambda_{(r, t-r)}^{(i, k-i)}=\sum_{(i, k-i) \in K \backslash R} u_{i} \cdot \lambda_{r}^{(i)} \cdot \bar{\lambda}_{t-r}^{(k-i)} .
$$

The number of blocks in $\mathfrak{B}_{(i, k-i)}^{*}$ containing $T_{(r, t-r)}$, for all $(i, k-i) \in R$, is then

$$
\sum_{(i, k-i) \in R} u_{i} . \Lambda_{(r, t-r)}^{*(i, k-i)}
$$

where

$$
\Lambda_{(r, t-r)}^{*(i, k-i)}=\left\{\begin{aligned}
\lambda_{r}^{*(i)} \cdot \bar{\lambda}_{t-r}^{*(k-i)} \cdot N_{i} \cdot z_{i} & \text { if } r \leq s_{i}, t-r \leq s_{k-i}, \\
\lambda_{r}^{*(i)} \cdot \bar{\lambda}_{t-r}^{(k-i)} \cdot z_{i} & \text { if } r \leq s_{i}, t-r>s_{k-i}, \\
\lambda_{r}^{(i)} \cdot \bar{\lambda}_{t-r}^{*(k-i)} \cdot z_{i} & \text { if } r>s_{i}, t-r \leq s_{k-i}
\end{aligned}\right.
$$

It follows that the number of blocks in $\mathfrak{B}$ containing $T_{r, t-r}$ is equal to

$$
L_{r, t-r}:=\sum_{(i, k-i) \in R} u_{i} \cdot \Lambda_{(r, t-r)}^{*(i, k-i)}+\sum_{(i, k-i) \in K \backslash R} u_{i} \cdot \Lambda_{(r, t-r)}^{(i, k-i)} .
$$

Since any $t$-subset of $X$ is of form $T_{r, t-r}$ for some $r \in\{0, \ldots, t\}$, we see that if

$$
L_{0, t}=L_{1, t-1}=\cdots=L_{t, 0}:=\Lambda
$$

for a positive integer $\Lambda$, then $(X, \mathfrak{B})$ forms a simple $t$-design with parameters $t$ $(v, k, \Lambda)$.

We record the construction above in the following theorem.
Theorem 3.1 Let $v, k, t$ be integers with $v>k>t \geq 2$. Let $X$ be a $v$-set and let $X=X_{1} \cup X_{2}$ be a partition of $X$ with $\left|X_{1}\right|=v_{1}$ and $\left|X_{2}\right|=v_{2}$. Let $D_{i}=\left(X_{1}, \mathfrak{B}^{(i)}\right)$ be the complete $i$ - $\left(v_{1}, i, 1\right)$ design for $i=0, \ldots, t$ and let $D_{i}=\left(X_{1}, \mathfrak{B}^{(i)}\right)$ be a simple $t-\left(v_{1}, i, \lambda_{t}^{(i)}\right)$ design for $i=t+1, \ldots, k$. Similarly, let $\bar{D}_{i}=\left(X_{2}, \overline{\mathfrak{B}}^{(i)}\right)$ be the complete $i$ - $\left(v_{2}, i, 1\right)$ design for $i=0, \ldots, t$, and let $\bar{D}_{i}=\left(X_{2}, \overline{\mathfrak{B}}^{(i)}\right)$ be a simple $t-\left(v_{2}, i, \bar{\lambda}_{t}^{(i)}\right)$ design for $i=t+1, \ldots, k$. Let $K=\{(0, k),(1, k-1), \ldots,(k-1,1),(k, 0)\}$. Suppose there exists a subset $R \subseteq K$ such that for each $(i, k-i) \in R$, the designs $D_{i}$ and $\bar{D}_{k-i}$ have an $s_{i}$-resolution with $N_{i}$ classes and an $s_{k-i}$-resolution with $\bar{N}_{k-i}$ classes, respectively, satisfying the following conditions.
(i) $N_{i}=\bar{N}_{k-i}$,
(ii) $s_{i}+s_{k-i} \geq 2\left\lfloor\frac{t}{2}\right\rfloor$.

## Define

$$
\mathfrak{B}=\bigcup_{(i, k-i) \in R} \mathfrak{B}_{(i, k-i)}^{*} \times\left[u_{i}\right] \cup \bigcup_{(i, k-i) \in K \backslash R} \mathfrak{B}_{(i, k-i)} \times\left[u_{i}\right],
$$

for $u_{i} \in\{0,1\}, i=0, \ldots, k$,
$\mathfrak{B}_{(i, k-i)}^{*}:=\left\{B_{i} \cup \bar{B}_{k-i} \mid B_{i} \in \mathfrak{A}_{h}^{(i)}, \bar{B}_{k-i} \in \overline{\mathfrak{A}}_{j}^{(k-i)}, \varepsilon_{i} \leq d\left(\mathfrak{A}_{h}^{(i)}, \mathfrak{A}_{j}^{(i)}\right) \leq w_{i}, \varepsilon_{i}=0,1 ; w_{i} \leq\left\lfloor\frac{N_{i}}{2}\right\rfloor\right\}$,
with $w_{i}$ as variable, where $\mathfrak{A}_{1}^{(i)}, \ldots, \mathfrak{A}_{N_{i}}^{(i)}$ are $s_{i}$-resolution classes of $D_{i}$, with $\left(X_{1}, \mathfrak{A}_{h}^{(i)}\right)$ as an $s_{i}-\left(v_{1}, i, \lambda_{s_{i}}^{*(i)}\right)$ design; and $\overline{\mathfrak{A}}_{1}^{(k-i)}, \ldots, \overline{\mathfrak{A}}_{N_{i}}^{(k-i)}$ are $s_{k-i}$-resolution classes of $\bar{D}_{k-i}$, with $\left(X_{2}, \overline{\mathfrak{A}}_{h}^{(k-i)}\right)$ as an $s_{k-i}-\left(v_{2}, k-i, \bar{\lambda}_{s_{k-i}}^{*(k-i)}\right)$ design; and

$$
\mathfrak{B}_{(i, k-i)}:=\left\{B=B_{i} \cup \bar{B}_{k-i} \mid B_{i} \in \mathfrak{B}^{(i)}, \bar{B}_{k-i} \in \overline{\mathfrak{B}}^{(k-i)}\right\} .
$$

Define

$$
L_{r, t-r}:=\sum_{(i, k-i) \in R} u_{i} \cdot \Lambda_{(r, t-r)}^{*(i, k-i)}+\sum_{(i, k-i) \in K \backslash R} u_{i} \cdot \Lambda_{(r, t-r)}^{(i, k-i)},
$$

for $r=0, \ldots, t$, where

$$
\Lambda_{(r, t-r)}^{*(i, k-i)}=\left\{\begin{aligned}
\lambda_{r}^{*(i)} \cdot \bar{\lambda}_{t-r}^{*(k-i)} \cdot N_{i} \cdot z_{i} & \text { if } r \leq s_{i}, t-r \leq s_{k-i} \\
\lambda_{r}^{*(i)} \cdot \bar{\lambda}_{t-r}^{(k-i)} \cdot z_{i} & \text { if } r \leq s_{i}, t-r>s_{k-i} \\
\lambda_{r}^{(i)} \cdot \bar{\lambda}_{t-r}^{*(k-i)} \cdot z_{i} & \text { if } r>s_{i}, t-r \leq s_{k-i}
\end{aligned}\right.
$$

with $z_{i}=\left(2 w_{i}+1-\varepsilon_{i}\right)$, if $w_{i}<\frac{N_{i}}{2}$, and $z_{i}=\left(2 w_{i}-\varepsilon_{i}\right)$, if $w_{i}=\frac{N_{i}}{2}$; and

$$
\Lambda_{(r, t-r)}^{(i, k-i)}=\lambda_{r}^{(i)} \cdot \bar{\lambda}_{t-r}^{(k-i)} .
$$

Assume that

$$
\begin{equation*}
L_{0, t}=L_{1, t-1}=\cdots=L_{t, 0}:=\Lambda \tag{3}
\end{equation*}
$$

for a positive integer $\Lambda$, then $(X, \mathfrak{B})$ is a simple $t-(v, k, \Lambda)$ design.
Remarks 3.1 1. In the basic construction the set $\mathfrak{B}_{(i, k-i)}$ of the new blocks is uniquely determined as the unions of all the pairs of blocks in $D_{i}$ and $\bar{D}_{k-i}$. Whereas in the construction using resolutions in Theorem 3.1 the set $\mathfrak{B}_{(i, k-i)}^{*}$ is no longer unique. Its size varies according to the variable $z_{i}$.
2. Theorem 3.1 does not restrict to constructing simple $t$-designs. Obviously, if any of the ingredient designs is non-simple, then the construction will yield non-simple designs.

## 4 Applications

In this section we illustrate the construction in Theorem 3.1 through a number of examples which show the strength of the method.

In the following we will employ the notation from Chapter II. $4: t$-Designs with $t \geq 3$ of the Handbook of Combinatorial Designs [12]. The parameter set $t-(v, k, \lambda)$ of a design will be written as $t-\left(v, k, m \lambda_{\min }\right)$. Since the supplement of a simple $t$ $(v, k, \lambda)$ design is a $t-\left(v, k, \lambda_{\max }-\lambda\right)$ design, we usually consider simple $t-(v, k, \lambda)$ designs with $\lambda \leq \lambda_{\max } / 2$. Thus, the upper limit of $m$ of a constructed design will be LIM $=\left\lfloor\lambda_{\text {max }} /\left(2 \lambda_{\text {min }}\right)\right\rfloor$. But, it should be remarked that, when an ingredient design with index $\lambda$ is used, then $\lambda$ can take on all possible values, i.e. $\lambda_{\min } \leq \lambda \leq \lambda_{\max }$.

### 4.1 Simple 5-(38, $k, \Lambda$ ) designs with $k=8,9,10$

We apply the construction in Theorem 3.1 to the cases $t=5, v_{1}=v_{2}=19$ and $k=8,9,10$.

### 4.1.1 Simple 5 -( $38,8, \Lambda$ ) designs

Here we show a detailed example to illustrate the construction.
Let $X=X_{1} \cup X_{2}$ be a partition of the point set $X$ with $|X|=38$ into two subsets $X_{1}$ and $X_{2}$ with $\left|X_{1}\right|=\left|X_{2}\right|=19$. For $i=0,1,2,3,4,5$ let $D_{i}=\left(X_{1}, \mathcal{B}^{(i)}\right)$ be the complete $i-\left(19, i, \lambda_{i}^{(i)}\right):=i-(19, i, 1)$ design. For $i=6,7,8$ let $D_{i}=\left(X_{1}, \mathcal{B}^{(i)}\right)$ be a simple 5-(19, $\left.i, \lambda_{5}^{(i)}\right)$ design. These designs have the following parameters.

- $5-\left(19,6, \lambda_{5}^{(6)}\right)=5-(19,6, m 2), m=1,2, \ldots, 7$.
- $5-\left(19,7, \lambda_{5}^{(7)}\right)=5-(19,7, m 7), m=1,2, \ldots, 13$
- $5-\left(19,8, \lambda_{5}^{(8)}\right)=5-(19,8, m 28), m=1,2, \ldots, 13$

Correspondingly, let $\bar{D}_{i}=\left(X_{2}, \overline{\mathcal{B}}^{(i)}\right)$ be simple designs defined on $X_{2}$. Here $K=$ $\{(0,8),(1,7),(2,6),(3,5),(4,4),(5,3),(6,2),(7,1),(8,0)\}$.

It is known that the complete designs $D_{i}$ and $\bar{D}_{i}$ for $i=3,4,5$ have each a 2-resolution with the number of resolution classes $N_{i}=17$, i.e. the large sets $\operatorname{LS}[17](2, i, 19)$, see for instance Chapter II. 4 [12]. We choose

$$
R=\{(3,5),(4,4),(5,3)\}
$$

Thus we have

- $\mathfrak{B}^{(3)}=\bigcup_{j=1}^{17} \mathfrak{A}_{j}^{(3)}$, where $\left(X_{1}, \mathfrak{A}_{j}^{(3)}\right)$ is a 2- $\left(19,3, \lambda_{2}^{*(3)}\right)=2-(19,3,1)$ design, and $\lambda_{3}^{(3)}=1, \lambda_{2}^{*(3)}=1, \lambda_{1}^{*(3)}=9, \lambda_{0}^{*(3)}=57$;
- $\mathfrak{B}^{(4)}=\bigcup_{j=1}^{17} \mathfrak{A}_{j}^{(4)}$, where $\left(X_{1}, \mathfrak{A}_{j}^{(4)}\right)$ is a $2-\left(19,4, \lambda_{2}^{*(4)}\right)=2-(19,4,8)$ design and $\lambda_{4}^{(4)}=1, \lambda_{3}^{(4)}=16, \lambda_{2}^{*(4)}=8, \lambda_{1}^{*(4)}=48, \lambda_{0}^{*(4)}=228$;
- $\mathfrak{B}^{(5)}=\bigcup_{j=1}^{17} \mathfrak{A}_{j}^{(5)}$, where $\left(X_{1}, \mathfrak{A}_{j}^{(i)}\right)$ is a $2-\left(19,5, \lambda_{2}^{*(4)}\right)=2-(19,5,40)$ design and $\lambda_{5}^{(5)}=1, \lambda_{4}^{(5)}=15, \lambda_{3}^{(5)}=120, \lambda_{2}^{*(5)}=40, \lambda_{1}^{*(5)}=180, \lambda_{0}^{*(5)}=684 ;$

Similarly, the complete designs $\bar{D}_{i}$ have the same 2-resolutions as $D_{i}$, each having $\bar{N}_{i}=17$ resolution classes, for $i=3,4,5$. Thus $\overline{\mathfrak{B}}^{(i)}=\bigcup_{j=1}^{17} \overline{\mathfrak{A}}_{j}^{(i)}$, and each $\left(X_{2}, \overline{\mathfrak{A}}_{j}^{(i)}\right)$ is a $2-\left(19, i, \bar{\lambda}_{2}^{*(i)}\right)$ design with $\bar{\lambda}_{2}^{*(i)}=\lambda_{2}^{*(i)}$.

We compute

$$
L_{r, 5-r}=\sum_{(i, 8-i) \in R} u_{i} \cdot \Lambda_{(r, 5-r)}^{*(i, 8-i)}+\sum_{(i, 8-i) \in K \backslash R} u_{i} \cdot \Lambda_{(r, 5-r)}^{(i, 8-i)},
$$

for $r=0, \ldots, 5$, and $u_{i}=0,1$. If $(i, 8-i) \in K \backslash R$, then

$$
\Lambda_{(r, 5-r)}^{(i, 8-i)}=\lambda_{r}^{(i)} \cdot \bar{\lambda}_{5-r}^{(8-i)} .
$$

If $(i, 8-i) \in R$, then the values of $\Lambda_{(r, 5-r)}^{*(i, 8-i)}$ are computed by using the formula

$$
\Lambda_{(r, t-r)}^{*(i, k-i)}=\left\{\begin{aligned}
\lambda_{r}^{*(i)} \cdot \bar{\lambda}_{t-r}^{*(k-i)} \cdot N_{i} \cdot z_{i} & \text { if } r \leq s_{i}, t \leq s_{k-i}, \\
\lambda_{r}^{* i)} \cdot \bar{\lambda}_{t-r}^{(k-i)} \cdot z_{i} & \text { if } r \leq s_{i}, t>s_{k-i}, \\
\lambda_{r}^{(i)} \cdot \bar{\lambda}_{t-r}^{*(k-i)} \cdot z_{i} & \text { if } r>s_{i}, t \leq s_{k-i} .
\end{aligned}\right.
$$

Here we have

$$
\begin{array}{lr}
\Lambda_{0,5}^{*(3,5)}=\lambda_{0}^{*(3)} \cdot \bar{\lambda}_{5}^{(5)} \cdot z_{3}=57 z_{3}, & \Lambda_{1,4}^{*(3,5)}=\lambda_{1}^{*(3)} \cdot \bar{\lambda}_{4}^{(5)} \cdot z_{3}=9 \times 15 z_{3}, \\
\Lambda_{2,3}^{*(3,5)}=\lambda_{2}^{*(3)} \cdot \bar{\lambda}_{3}^{(5)} \cdot z_{3}=120 z_{3}, & \Lambda_{3,2}^{*(3,5)}=\lambda_{3}^{(3)} \cdot \bar{\lambda}_{2}^{(5)} \cdot z_{3}=40 z_{3}, \\
\Lambda_{4,1}^{*(3,5)}=\Lambda_{5,0}^{*(3,5)}=0 . &
\end{array}
$$

$$
\begin{array}{lr}
\Lambda_{0,5}^{*(5,3)}=\Lambda_{1,4}^{*(5,3)}=0, & \Lambda_{2,3}^{*(5,3)}=\lambda_{2}^{*(5)} \cdot \bar{\lambda}_{3}^{(3)} \cdot z_{5}=40 z_{5}, \\
\Lambda_{3,2}^{*(5,3)}=\lambda_{3}^{(5)} \cdot \bar{\lambda}_{2}^{*(3)} \cdot z_{5}=120 z_{5}, & \Lambda_{4,1}^{*(5,3)}=\lambda_{4}^{(5)} \cdot \bar{\lambda}_{1}^{*(3)} \cdot z_{5}=15 \times 9 z_{5}, \\
\Lambda_{5,0}^{*(5,3)}=\lambda_{5}^{(5)} \cdot \bar{\lambda}_{0}^{*(3)} \cdot z_{5}=57 z_{5} . & \\
\Lambda_{0,5}^{*(4,4)}=\Lambda_{5,0}^{*(4,4)}=0, & \Lambda_{1,4}^{*(4,4)}=\lambda_{1}^{*(4)} \cdot \bar{\lambda}_{4}^{(4)} \cdot z_{4}=48 z_{4}, \\
\Lambda_{2,3}^{*(4,4)}=\lambda_{2}^{*(4)} \cdot \bar{\lambda}_{3}^{(4)} \cdot z_{4}=8 \times 16 z_{4}, & \Lambda_{3,2}^{*(4,4)}=\lambda_{3}^{(4)} \cdot \bar{\lambda}_{2}^{*(4)} \cdot z_{4}=16 \times 8 z_{4}, \\
\Lambda_{4,1}^{*(4,4)}=\lambda_{4}^{(4)} \cdot \bar{\lambda}_{1}^{*(4)} \cdot z_{4}=48 z_{4} . &
\end{array}
$$

It follows that

$$
\begin{aligned}
& L_{0,5}=u_{0} \bar{\lambda}_{5}^{(8)}+u_{1} 19 \bar{\lambda}_{5}^{(7)}+u_{2} 171 \bar{\lambda}_{5}^{(6)}+u_{3} 57 z_{3} \\
& L_{1,4}=u_{1} 5 \bar{\lambda}_{5}^{(7)}+u_{2} 9 \times 15 \bar{\lambda}_{5}^{(6)}+u_{3} 9 \times 15 z_{3}+u_{4} 48 z 4 \\
& L_{2,3}=u_{2} 40 \bar{\lambda}_{5}^{(6)}+u_{3} 120 z_{3}+u_{4} 8 \times 16 z_{4}+u_{5} 40 z 5 \\
& L_{3,2}=u_{6} 40 \lambda_{5}^{(6)}+u_{5} 120 z_{5}+u_{4} 16 \times 8 z 4+u_{3} 40 z_{3} \\
& L_{4,1}=u_{7} 5 \lambda_{5}^{(7)}+u_{6} 15 \times 9 \lambda_{5}^{(6)}+u_{5} 15 \times 9 z_{5}+u_{4} 48 z 4 \\
& L_{5,0}=u_{8} \lambda_{5}^{(8)}+u_{7} 19 \lambda_{5}^{(7)}+u_{6} 171 \lambda_{5}^{(6)}+u_{5} 57 z_{5}
\end{aligned}
$$

Each set of values of $u_{i} \in\{0,1\}, i=0, \ldots, 8 ; z 3, z 4, z 5=1, \ldots, 17 ; \lambda_{5}^{(j)}$ and $\bar{\lambda}_{5}^{(j)}$, $j=6,7,8$ for which the equalities

$$
L_{0,5}=L_{1,4}=L_{2,3}=L_{3,2}=L_{4,1}=L_{5,0}:=\Lambda
$$

is satisfied for a positive integer $\Lambda$ will yield a simple $5-(38,8, \Lambda)$ design. Recall that a $5-(38,8, \Lambda)$ can be written as $5-(38,8, m 4)$ with $\lambda_{\min }=4$ and $\lambda_{\max }=5456$. Thus LIM $=\lfloor 5456 / 2 * 4\rfloor=682$. By solving the equalities above we obtain all solutions for $m \leq 1364$. Altogether 33 values for $m$ have been found, of which 16 values of $m \leq$ LIM. Since, not all simple $5-\left(19, i, \lambda_{5}^{(i)}\right)$ designs are known to exist, for example, $5-(19,7, m 7)$ designs are known for $m=4,5,6,7,8,9,13$ only, we just obtain the following 5 new simple 5 - $(38,8, m 4)$ designs for $m=280,488,524,560,560$ (the number 560 repeats twice, as we have two distinct non isomorphic solutions for this value of $m$ ). The details of these 5 constructed designs are given in Table 1.

Table 1: Constructed simple 5- $(38,8, \Lambda)$ designs

| $m$ | $z_{3}$ | $z_{4}$ | $z_{5}$ | $\lambda_{5}^{(6)}$ | $\lambda_{5}^{(7)}$ | $\lambda_{5}^{(8)}$ |
| :--- | ---: | ---: | ---: | ---: | :---: | ---: |
| 280 | 7 | 0 | 7 | 0 | 35 | 56 |
| 488 | 8 | 4 | 8 | 4 | 28 | 280 |
| 524 | 6 | 7 | 6 | 6 | 28 | 196 |
| 560 | 4 | 10 | 4 | 8 | 28 | 112 |
| 560 | 9 | 5 | 9 | 4 | 49 | 112 |

An entry 0 in a column of the table implies that $u_{i}=0$, otherwise $u_{i}=1$. Here we have $\lambda_{5}^{(j)}=\bar{\lambda}_{5}^{(j)}, j=6,7,8$ for all these solutions.

### 4.1.2 Simple 5 -( $38, k, \Lambda$ ) designs with $k=9,10$

Again we assume that $v_{1}=v_{2}=19$ for the construction of simple 5 - $(38, k, \Lambda)$ designs with $k=9,10$.

- For construction of $5-(38,9, \Lambda)=5-(38,9, m 30)$ designs with LIM $=682$, we make use of the large sets $\operatorname{LS}[17](2, i, 19), i=3,4,5,6$, i.e. the 2 -resolutions of the complete designs $i-(19, i, 1)$ with resolution class number $N_{i}=17$. Thus, we have $R=\{(3,6),(4,5),(5,4),(6,3)\}$. And the equalities $L_{r, t-r}$ are the following.

$$
\begin{aligned}
L_{0,5} & =u_{0} \bar{\lambda}_{5}^{(9)}+u_{1} 19 \bar{\lambda}_{5}^{(8)}+u_{2} 171 \bar{\lambda}_{5}^{(7)}+u_{3} 57 \times 14 z_{3}+u_{4} 228 z_{4}, \\
L_{1,4} & =u_{1} 15 \bar{\lambda}_{5}^{(8)} / 4+u_{2} 18 \times 5 \bar{\lambda}_{5}^{(7)}+u_{3} 9 \times 105 z_{3}+u_{4} 48 \times 15 z_{4}+u_{5} 180 z_{5}, \\
L_{2,3} & =u_{2} 20 \bar{\lambda}_{5}^{(7)}+u_{3} 560 z_{3}+u_{4} 8 \times 120 z_{4}+u_{5} 40 \times 16 z_{5}+u_{6} 140 z_{6}, \\
L_{3,2} & =u_{7} 20 \lambda_{5}^{(7)}+u_{6} 560 z_{6}+u_{5} 120 \times 8 z_{5}+u_{4} 16 \times 40 z_{4}+u_{3} 140 z_{3}, \\
L_{4,1} & =u_{8} 15 \lambda_{5}^{(8)} / 4+u_{7} 5 \times 18 \lambda_{5}^{(7)}+u_{6} 105 \times 9 z_{6}+u_{5} 15 \times 48 z_{5}+u_{4} 180 z_{4}, \\
L_{5,0} & =u_{9} \lambda_{5}^{(9)}+u_{8} 19 \lambda_{5}^{(8)}+u_{7} 171 \lambda_{5}^{(7)}+u_{6} 14 \times 57 z_{6}+u_{5} 228 z_{5} .
\end{aligned}
$$

Solving the equalities $L_{0,5}=L_{1,4}=L_{2,3}=L_{3,2}=L_{4,1}=L_{5,0}=\Lambda$ for $\Lambda>0$ with respect to $z_{i}=1, \ldots, 17$ we obtain 20 values for $m$ with $m \leq$ LIM leading to simple $5-(38,9, \Lambda)=5-(38,9, m 30)$ designs. Of which 14 designs can be constructed whose details are given in Table 2.

Table 2: Constructed simple 5- $(38,9, \Lambda)$ designs

| $m$ | $z_{3}$ | $z_{4}$ | $z_{5}$ | $z_{6}$ | $\lambda_{5}^{(7)}$ | $\lambda_{5}^{(8)}$ | $\lambda_{5}^{(9)}$ | $\bar{\lambda}_{5}^{(7)}$ | $\bar{\lambda}_{5}^{(8)}$ | $\bar{\lambda}_{5}^{(9)}$ |
| :--- | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: |
| 100 | 2 | 1 | 1 | 2 | 0 | 56 | 112 | 0 | 56 | 112 |
| 200 | 4 | 2 | 2 | 4 | 0 | 112 | 224 | 0 | 112 | 224 |
| 300 | 6 | 3 | 3 | 6 | 0 | 168 | 336 | 0 | 168 | 336 |
| 400 | 8 | 4 | 4 | 8 | 0 | 224 | 448 | 0 | 224 | 448 |
| 402 | 5 | 5 | 5 | 5 | 28 | 84 | 546 | 28 | 84 | 546 |
| 500 | 10 | 5 | 5 | 10 | 0 | 280 | 560 | 0 | 280 | 560 |
| 502 | 7 | 6 | 6 | 7 | 28 | 140 | 658 | 28 | 140 | 658 |
| 504 | 4 | 7 | 7 | 4 | 56 | 0 | 756 | 56 | 0 | 756 |
| 582 | 10 | 4 | 11 | 3 | 28 | 168 | 588 | 63 | 84 | 189 |
| 602 | 9 | 7 | 7 | 9 | 28 | 196 | 770 | 28 | 196 | 770 |
| 604 | 6 | 8 | 8 | 6 | 56 | 56 | 868 | 56 | 56 | 868 |
| 660 | 9 | 8 | 8 | 9 | 35 | 252 | 21 | 35 | 252 | 21 |
| 680 | 8 | 11 | 4 | 15 | 0 | 364 | 602 | 35 | 280 | 203 |
| 682 | 5 | 12 | 5 | 12 | 28 | 224 | 700 | 63 | 140 | 301 |

It should be noted that when applying the basic construction for $t=5, v_{1}=$ $v_{2}=19$ and $k=8,9$ we only obtain the trivial solutions, namely the complete
$5-(38,8,1364 \times 4)$ and $5-(38,9,1364 \times 30)$ designs. This could be explained as follows.

In general, if $k \leq 2 t-1$, then one of the designs in each pair $\left(D_{i}, \bar{D}_{k-i}\right)$ is either the empty or the trivial design and at least one pair having both the trivial designs, therefore it leaves little room for the basic construction to produce a non-trivial solution, unless many pairs are unused, i.e. $u_{i}=0$. The construction using resolutions indeed makes more room to create non-trivial solutions, as we have seen in the above examples.

- For construction of $5-(38,10, \Lambda)=5-(38,10, m 6)$ designs with $\operatorname{LIM}=19778$, we again employ the 2 -resolutions of the complete designs $i-(19, i, 1)$ for $i=$ $3,4,5,6,7$ with resolution class number $N_{i}=17$. Here,

$$
R=\{(3,7),(4,6),(5,5),(6,4),(7,3)\}
$$

And we have

$$
\begin{aligned}
L_{0,5}= & u_{0} \bar{\lambda}_{5}^{(10)}+u_{1} 19 \bar{\lambda}_{5}^{(9)}+u_{2} 171 \bar{\lambda}_{5}^{(8)}+u_{3} 57 \times 91 z_{3}+u_{4} 228 \times 14 z_{4}+u_{5} 684 z_{5} \\
L_{1,4}= & u_{1} 3 \bar{\lambda}_{5}^{(9)}+u_{2} 18 \times 15 \bar{\lambda}_{5}^{(8)} / 4+u_{3} 9 \times 455 z_{3}+u_{4} 48 \times 105 z_{4}+u_{5} 180 \times 15 z_{5} \\
& +u_{6} 504 z_{6} \\
L_{2,3}= & u_{2} 12 \bar{\lambda}_{5}^{(8)}+u_{3} 1820 z_{3}+u_{4} 8 \times 560 z_{4}+u_{5} 40 \times 120 z_{5}+u_{6} 140 \times 16 z_{6} \\
& +u_{7} 364 z_{7}, \\
L_{3,2}= & u_{8} 12 \lambda_{5}^{(8)}+u_{7} 1820 z_{7}+u_{6} 560 \times 8 z_{6}+u_{5} 120 \times 40 z_{5}+u_{4} 16 \times 140 z_{4} \\
& +u_{3} 364 z_{3} \\
L_{4,1}= & u_{9} 3 \lambda_{5}^{(9)}+u_{8} 15 \times 18 \lambda_{5}^{(8)} / 4+u_{7} 455 \times 9 z_{7}+u_{6} 105 \times 48 z_{6}+u_{5} 15 \times 180 z_{5} \\
& +u_{4} 504 z_{4}, \\
L_{5,0}= & u_{10} \lambda_{5}^{(10)}+u_{9} 19 \lambda_{5}^{(9)}+u_{8} 171 \lambda_{5}^{(8)}+u_{7} 91 \times 57 z_{7}+u_{6} 14 \times 228 z_{6}+u_{5} 684 z_{5} .
\end{aligned}
$$

Solving the equalities $L_{0,5}=L_{1,4}=L_{2,3}=L_{3,2}=L_{4,1}=L_{5,0}=\Lambda$ for $\Lambda>0$ with respect to $z_{i}=1, \ldots, 17$ we obtain an entire number of 479 solutions, of which 239 have $m \leq$ LIM. From these 239 parameters 131 simple $5-(38,10, m 6)$ designs have been shown to exist. The values of $m$ for these designs are

| 12768 | 17416 | 2604 | 6076 | 7252 | 10724 | 13668 | 15108 |
| ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: |
| 15372 | 18580 | 18844 | 3768 | 6976 | 8416 | 8680 | 11624 |
| 11888 | 12152 | 16272 | 16536 | 16800 | 19744 | 4932 | 8404 |
| 9580 | 9844 | 12788 | 13052 | 13316 | 13580 | 17172 | 17436 |
| 17700 | 17964 | 18228 | 6096 | 9040 | 10480 | 10744 | 11008 |
| 13952 | 11536 | 14216 | 14480 | 15920 | 18600 | 18864 | 19128 |
| 7260 | 11644 | 11908 | 12172 | 14852 | 15116 | 15380 | 15644 |
| 16556 | 16820 | 19500 | 17084 | 19764 | 8424 | 11368 | 12544 |
| 12808 | 13072 | 13336 | 16016 | 16280 | 16544 | 16808 | 17984 |
| 9060 | 9588 | 13972 | 14236 | 16916 | 14500 | 17180 | 17444 |
| 17708 | 18884 | 19148 | 10224 | 10752 | 14872 | 15136 | 15400 |
| 18080 | 15664 | 18344 | 18608 | 19520 | 11388 | 11916 | 16036 |
| 16300 | 16564 | 19244 | 16828 | 19508 | 19772 | 12552 | 13080 |
| 17200 | 17464 | 17728 | 13716 | 14244 | 18100 | 18364 | 18628 |
| 18892 | 19156 | 14880 | 15408 | 19264 | 19528 | 16044 | 17208 |
| 18384 | 18372 | 19536 | 16844 | 11316 | 13908 | 14280 | 14808 |
| 19720 | 16872 | 17772 |  |  |  |  |  |

Here are two examples:

- $5-(38,10,2604 \times 6)$ with $z_{3}=1, z_{4}=2, z_{6}=2, z_{7}=1, \bar{\lambda}_{5}^{(9)}=147, \bar{\lambda}_{5}^{(10)}=1260$, $u_{2}=u_{5}=u_{8}=0$ and $\lambda_{5}^{(i)}=\bar{\lambda}_{5}^{(i)}$ for $i=9,10$.
- 5- $(38,10,11316 \times 6)$ with $z_{3}=2, z_{4}=8, z_{5}=2, z_{6}=7, z_{7}=4, \bar{\lambda}_{5}^{(8)}=140$, $\bar{\lambda}_{5}^{(9)}=336, \bar{\lambda}_{5}^{(10)}=294, \lambda_{5}^{(8)}=84, \lambda_{5}^{(9)}=378, \lambda_{5}^{(10)}=1890$.

On the other hand, when the basic construction is applied for this case (i.e. $v_{1}=$ $v_{2}=19$ and $k=10$ ), we just obtain 5 solutions with $m \leq$ LIM.

Remark 4.1 1. It should be noted that when $v_{1}=v_{2}$, any solution with $\lambda_{t}^{(i)} \neq \bar{\lambda}_{t}^{(i)}$ will appear twice by reason of symmetry, since $\lambda_{t}^{(i)}$ and $\bar{\lambda}_{t}^{(i)}$ may be interchanged. These two solutions are indeed the same. This fact should be taken into account by counting the number of solutions throughout Section 4.
2. Up to now the number of known parameter sets for $5-(38, k, \Lambda)$ with $k=8,9,10$ are 8,14 , and 23 respectively, see [12], for instance. For $k=8,9$ all the parameters of the constructed designs differ from the known ones. For $k=10$, only one of the 23 known parameter sets does appear in the list of 131 constructed designs, namely the parameters $5-(38,10,11368 \times 6)$. However, it is not known whether the corresponding designs are isomorphic.

### 4.2 Some further results of applications

We briefly record some further examples of simple $t$-designs for $t=4,5,6$ by using Theorem 3.1.

### 4.2.1 $\quad t=4$

Following are several small parameters for $t=4$.

1. $4-(26,8, m 35)$ : Take $v_{1}=v_{2}=13$ and $R=\{(3,5),(4,4),(5,3)\}$ by using $\operatorname{LS}[55](2,4,13)$ and $\operatorname{LS}[11](2, i, 13)$ for $i=3,5$. There are 3 non-trivial solutions of $\mathrm{Eq}(3)$ with $m=44,66$ satisfying $m \leq \operatorname{LIM}(=104)$. A design with $m=44$ is known. The two solutions for $m=66$ are non-isomorphic and new. These are

- $u_{4}=0, z_{3}=z_{5}=7, \lambda_{4}^{(7)}=42, \lambda_{4}^{(8)}=126, u_{2}=u_{6}=0$, and $\bar{\lambda}_{4}^{(i)}=\lambda_{4}^{(i)}$ for $i=7,8$.
- $z_{4}=24, z_{3}=z_{5}=2, \lambda_{4}^{(6)}=18, \lambda_{4}^{(8)}=126, u_{1}=u_{7}=0$, and $\bar{\lambda}_{4}^{(i)}=\lambda_{4}^{(i)}$ for $i=6,8$.

The basic construction for $4-(26,8, m 35)$ with $v_{1}=v_{2}=13$ only yields the trivial solution.
2. 4-(28, $9, m 168)$ : Take $v_{1}=v_{2}=14$ and $R=\{(4,5),(5,4)\}$ by using $\operatorname{LS}[11](2, i, 14)$ for $i=4,5$. There is a unique non-trivial solution of $\mathrm{Eq}(3)$ with $m=110$ satisfying $m \leq \operatorname{LIM}(=126)$. This solution with $z_{4}=z_{5}=4, u_{2}=u_{7}=0, \lambda_{4}^{(6)}=30$, $\lambda_{4}^{(8)}=210, \lambda_{4}^{(9)}=252$, and $\bar{\lambda}_{4}^{(i)}=\lambda_{4}^{(i)}$ for $i=6,8,9$ yields a new design.
3. $4-(30,7, m 20)$ : Take $v_{1}=v_{2}=15$ and $R=\{(3,4),(4,3)\}$ by using $\operatorname{LS}[13](2, i, 15)$ for $i=3,4$. There are 3 non-trivial solutions of $\operatorname{Eq}(3)$ with $m=39,52,65$ satisfying $m \leq \operatorname{LIM}(=65)$. The solution for $m=52$ with $z_{3}=z_{5}=5, \lambda_{4}^{(5)}=5$, $\lambda_{4}^{(6)}=15, \lambda_{4}^{(7)}=115$, and $\bar{\lambda}_{4}^{(i)}=\lambda_{4}^{(i)}$ for $i=5,6,7$ gives a new design.

### 4.2.2 $t=5$

1. $5-(36,10, m 63)$ : Take $v_{1}=v_{2}=18$ and $R=\{(5,5)\}$ by using $\operatorname{LS}[7](2,5,18)$. There are 164 non-trivial solutions of $\operatorname{Eq}(3)$ with $m \leq \operatorname{LIM}(=1348)$. Of which 37 are shown to exist. It is interesting to remark that these 37 designs include the 10 designs constructed using the basic construction [39]. Actually, 27 new designs with parameters $5-(36,10, m 63)$ have been obtained. These are

$$
\begin{aligned}
m= & 611,818,921,945,969,1048,1072,911,934,1094,1197,1221,1245,1269 \\
& 1324,1325,1348,1187,1210,1234,1337,1152,1176,1200,1224,1303,1131 .
\end{aligned}
$$

2. 5-(37, $8, m 40)$ : Take $v_{1}=13, v_{2}=24$ and $R=\{(3,5),(4,4),(5,3)\}$ by using $\operatorname{LS}[11](2, i, 13), \operatorname{LS}[11](2, i, 24)$ for $i=3,4,5$. There is a unique non-trivial solution of $\mathrm{Eq}(3)$ with $m=55$ such that $m \leq \operatorname{LIM}(=62)$. This solution with $z_{3}=2, z_{5}=2, z_{4}=8, \lambda_{5}^{(6)}=4, \lambda_{5}^{(7)}=28, \lambda_{5}^{(8)}=56, \bar{\lambda}_{5}^{(6)}=13, \bar{\lambda}_{5}^{(7)}=36$, $\bar{\lambda}_{5}^{(7)}=666$ gives a new design.
3. $5-(37,9, m 10):$ Take $v_{1}=13, v_{2}=24$ and $R=\{(3,6),(4,5),(5,4),(6,3)\}$ by using $\operatorname{LS}[11](2, i, 13), \operatorname{LS}[11](2, i, 24)$ for $i=3,4,5,6$. There is a unique nontrivial solution of $\operatorname{Eq}(3)$ with $m=874$ such that $m \leq \operatorname{LIM}(=1798)$. This solution with $z_{3}=2, z_{4}=2, z_{5}=4, z_{6}=1, \lambda_{5}^{(7)}=14, u_{8}=u_{9}=0, \bar{\lambda}_{5}^{(6)}=72$, $\bar{\lambda}_{5}^{(7)}=30, \bar{\lambda}_{5}^{(8)}=1980$ gives a new design.
4. 5-(44, $8, m)$ : Take $v_{1}=v_{2}=22$ and $R=\{(4,4)\}$ by using $\operatorname{LS}[19](2,5,22)$. There are 9 non-trivial solutions of $\mathrm{Eq}(3)$ with $m \leq \operatorname{LIM}(=4569)$. Of which one design with $m=3344$ and $u_{3}=u_{5}=0, z_{4}=4, \lambda_{5}^{(6)}=12, \lambda_{5}^{(7)}=16$, $\lambda_{5}^{(8)}=220$, and $\bar{\lambda}_{5}^{(i)}=\lambda_{5}^{(i)}$ for $i=6,7,8$, is shown to exist.
5. 5-(46, 10, m2): Take $v_{1}=v_{2}=23$ and $R=\{(4,6),(5,5),(6,4)\}$ by using $\operatorname{LS}[133](2,5,23), \mathrm{LS}[7](2, i, 23)$ for $i=4,6$. There are 3986 non-trivial solutions of $\operatorname{Eq}(3)$ with $m \leq \operatorname{LIM}(=187349)$. Of which 176 designs are shown to exist with the following values of $m$.

| 65246 | 75487 | 73758 | 83999 | 86526 | 94240 | 96140 | 96767 |
| ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: |
| 106381 | 107008 | 116622 | 117021 | 125134 | 123405 | 127262 | 139004 |
| 137503 | 142633 | 139403 | 143887 | 149644 | 151772 | 159885 | 162013 |
| 164540 | 166668 | 174781 | 185497 | 59014 | 79667 | 78166 | 89908 |
| 88179 | 92435 | 98420 | 102676 | 108661 | 110561 | 120802 | 126160 |
| 122930 | 125058 | 131043 | 136401 | 133171 | 139555 | 137826 | 145540 |
| 153425 | 157054 | 158308 | 175807 | 174705 | 184319 | 182818 | 184946 |
| 77064 | 94088 | 99446 | 96216 | 104329 | 109687 | 102600 | 119928 |
| 112841 | 116698 | 121828 | 115368 | 123082 | 132069 | 124982 | 125609 |
| 130967 | 142310 | 135223 | 140581 | 135850 | 141208 | 145464 | 150822 |
| 145863 | 153976 | 156104 | 163590 | 167846 | 166345 | 171475 | 168245 |
| 178486 | 183844 | 180614 | 185972 | 86526 | 105355 | 113867 | 112366 |
| 118750 | 124108 | 126635 | 134349 | 127262 | 136249 | 134748 | 131518 |
| 136876 | 137503 | 146490 | 139403 | 147117 | 140030 | 156731 | 149644 |
| 155002 | 151772 | 157130 | 150271 | 153900 | 159885 | 165243 | 162013 |
| 167371 | 166668 | 177612 | 174382 | 179512 | 182267 | 185896 | 187150 |
| 121505 | 119776 | 128288 | 132544 | 138529 | 142785 | 148770 | 145540 |
| 150670 | 160911 | 163039 | 159809 | 171152 | 169423 | 164692 | 170050 |
| 168321 | 179664 | 174705 | 182818 | 184946 | 135850 | 148694 | 152950 |
| 158707 | 163191 | 165091 | 163590 | 165718 | 171076 | 175332 | 179189 |
| 175959 | 181317 | 174230 | 185573 | 178486 | 183844 | 180614 | 185972 |
| 184471 | 182742 | 176985 | 164540 | 177612 | 179512 | 187226 | 173052 |

Here is an example with $m=59014: z_{4}=z_{6}=1, z_{5}=20, u_{2}=u_{8}=0$, $\lambda_{5}^{(7)}=36, \lambda_{5}^{(9)}=810, \lambda_{5}^{(10)}=7812$, and $\bar{\lambda}_{5}^{(i)}=\lambda_{5}^{(i)}$ for $i=7,9,10$.

### 4.2.3 $\quad t=6$

Following are some examples for $t=6$.

1. $6-(38,10, m 10)$ : Take $v_{1}=v_{2}=19$ and $R=\{(4,6),(5,5),(6,4)\}$ by using $\operatorname{LS}[4](3, i, 19)$ for $i=4,5,6$. There are 4 non-trivial solutions of $\mathrm{Eq}(3)$ with $m=1360,892,1340,1788$ for $m \leq \operatorname{LIM}(=1798)$.
2. 6-(46, 12, $m 420$ ): Take $v_{1}=v_{2}=23$ and $R=\{(6,6)\}$ by using $\operatorname{LS}[3](3,6,23)$. There are 2 non-trivial solutions of $\operatorname{Eq}(3)$ with $m=3363$, 3819 for $m \leq \operatorname{LIM(=}$ 4569). The solution for $m=3363$ has $z_{6}=1, \lambda_{5}^{(7)}=7, \lambda_{5}^{(8)}=40, \lambda_{5}^{(9)}=340$, $\lambda_{5}^{(10)}=350, \lambda_{5}^{(11)}=4046, \lambda_{5}^{(12)}=5320$, and $\bar{\lambda}_{5}^{(i)}=\lambda_{5}^{(i)}$ for $i=7,8,9,10,11,12$. All the ingredient designs corresponding to $m=3363$ exist except that the existence of a $6-(23,10,5 \times 70)$ design is still in doubt. So, we would have a $6-(46,12,3364 \times 420)$ design if a $6-(23,10,5 \times 70)$ design would exist.
3. $6-(50,12, m 308)$ : Take $v_{1}=v_{2}=25$ and $R=\{(6,6)\}$ by using $\operatorname{LS}[7](3,6,25)$. There are 195 non-trivial solutions of $\operatorname{Eq}(3)$ for $m \leq \operatorname{LIM(=11459).~}$

## 5 Conclusion

We have presented a recursive construction for simple $t$-designs by using the concept of resolutions. This may be viewed as an extension of the basic construction as shown in our previous paper. The $s$-resolutions of trivial $t$-designs are equivalent to the large sets of $s$-designs, which have been extensively studied. Since our construction does not exclude the use of trivial designs as ingredients, we have restricted its applications to resolutions of the trivial ingredient designs only. In spite of this fact, the construction still produces a large number of new simple $t$-designs. We strongly believe that the construction would unfold its full impact when we would gain more knowledge about resolutions of non-trivial $t$-designs.

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